

# Evolution of Cooperative Problem-Solving in an Artificial Economy

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## Abstract

We address the problem of how to reinforcement learn in ultra-complex environments, with huge state spaces, where one must learn to exploit compact structure of the problem domain. The approach we propose is to simulate the evolution of an artificial economy of computer programs. The economy is constructed based on two simple principles so as to assign credit to the individual programs for collaborating on problem solutions. We find empirically that, starting from programs that are random computer code, we are able to evolve systems that solve hard problems. In particular our economy has learned to solve almost all random Blocks World problems with goal stacks 200 blocks high. Competing methods solve such problems only up to goal stacks of at most 8 blocks. Our economy has also learned to unscramble about half a randomly scrambled Rubik's cube, and to solve several among a collection of commercially sold puzzles.

Keywords: Reinforcement learning, Artificial Economy, Artificial Life, Hayek, Tierra, S-expression, planning, search, macro-actions, emergent computation, evolutionary economics, code evolution, autonomous programming, Blocks World, Rubik's Cube, Rush Hour, evolutionary programming, classifier systems, genetic programming.

## 1 Introduction

Reinforcement learning(c.f. [SB98]) formalizes the problem of learning from interaction with an environment. The two standard approaches to RL are 'value iteration' and 'policy iteration'. At the current state of the art, neither appears to offer much hope for addressing ultra-complex problems. Standard approaches to value iteration depend on enumerating the state space, and are thus problematic when the state space is huge. Moreover, they depend on finding an evaluation function showing progress after a single action, and such a function may be extremely hard to learn, or even to represent. We believe that what is necessary in hard domains is to learn and exploit compact structure of the state space, for which powerful methods are yet to be discovered. Policy iteration, on the other hand, suffers from the fact that the space of policies is also enormous, and also has a bumpy fitness landscape. Standard methods for policy iteration, among which we would include the various flavors of evolutionary and genetic programming, thus grind to a halt on most complex problems.

The approach we propose here is to evolve an artificial economy of modules. The artificial economy is constructed so as to assign credit to the individual modules for their contribution.

Evolution then can progress by finding effective modules. This divide and conquer approach greatly expedites the search of program (a.k.a. policy) space and can result in a compact solution exploiting the structure of the problem.

Holland’s seminal Classifier Systems [Hol86] previously pursued a similar approach, using an economic model to assign credit to modules. Unfortunately, Classifier Systems have never succeeded in their goal of dynamically chaining modules together to solve interesting problems [WG89]. We believe we have understood the problems with Holland’s proposal, and other multiagent proposals, and how to correct them in terms of simple but fundamental principles. Our picture applies widely to the evolution of multi-agent systems, including ecologies and economies. We discuss here evidence for our picture, evolving artificial economies in three different representations and for three problems, that dynamically chain huge sequences of learned modules to solve problems vastly too large for competing methods to address. Control experiments show that if our system is modified to violate our simple principles, even in ways suggested as desirable in the Classifier System literature, performance immediately breaks as we would predict.

We present results here on three problems: Blocks World, Rubik’s Cube, and a commercially sold set of puzzles called “Rush Hour”. Block’s World has a lengthy history as a benchmark problem which we survey in section 3. Methods like TD learning using a neural net evaluation, genetic programming, inductive logic programming, and others are able to solve problems directly comparable to the ones we experiment with only when they are very small— containing a handful of blocks, i.e. before the exponential growth of the state space bites in. By contrast, we report solution of problems with hundreds of blocks, and state spaces containing of order  $10^{100}$  states, with a single reward state. On Rubik’s cube we have been able to evolve an economy able to unscramble about half of a randomly scrambled cube— comparable to the performance of many humans. On Rush Hour our approach has solved 15 different commercially sold puzzles, including one dubbed “advanced” by the manufacturer. Genetic Programming was unable to make headway on either Rubik or Rush Hour in our experiments, and we know of no other learning approach capable of interesting results here.

Note that we are attempting to learn a program capable of solving a class of problems. For example, we are attempting to learn a program capable of unscrambling a randomly scrambled Rubik’s cube. We conjecture that for truly complex problems, learning to solve a class is often the most effective way to solve an individual problem. You can’t learn to unscramble a Rubik’s cube, for example, without learning how to unscramble any Rubik’s cube. Likewise, the planning approach to Blocks World has attempted huge search to solve individual Blocks World problems, but what is needed is to learn to exploit the structure of Blocks World. In our comparison experiments, we have attempted to use alternative methods such as GP to learn to solve a class of problems, but failed. We train by presenting small instances, and our method learns to exploit the underlying compact structure of the problem domain. We conjecture that learning to exploit underlying compact structure is closely related to “understanding”, a human ability little explained.

Section 2 describes our artificial economy and the principles that it embodies. Section 3 describes the Blocks World problem and surveys alternative approaches. Section 4 describes the syntax of our programs. Section 5 describes our results on Blocks World. Section 6 describes a more complex economic model involving meta-learning. Section 7 describes our results with Rubik’s Cube. Section 8 describes our results with the RushHour Problem. Section 9 sums up what we have learned from these experiments and suggests avenues for further exploration. Appendix A discusses results on Blocks World with alternative methods such as TD learning and Genetic Programming. Appendix B gives pseudocode. Appendix C discusses parameter settings.

## 2 The Artificial Economy

This section will describe and motivate our artificial economy, which we call Hayek. Hayek interacts with a world that it may sense, where it may take actions, and that makes a payoff when it is put in an appropriate state. Blocks World, c.f. figure 1, is an example world. Hayek consists of a collection of modules, each consisting of a computer program with an associated numeric “wealth”. We call the modules “agents”<sup>1</sup> because we allow them to sense features of the world, compute, and take actions on the world. The system acts in a series of *auctions*. In each auction, each agent simulates the execution of its program on the current world and returns a non-negative number. This number can be thought of as the agent’s estimate of the value of the state its execution would reach. The agent bids an amount equal to the minimum of its wealth and the returned number. The solver with highest bid wins the auction. It pays this bid to the winner of the previous auction, executes its actions on the current world, and collects any reward paid by the world as well as the winning bid in the subsequent auction. Evolutionary pressure pushes agents to reach highly valued states and to bid accurately, lest they be outbid.

In each auction, each agent that has wealth more than a fixed sum  $10W_{\text{init}}$  creates a new agent that is a mutation of itself. Like an investor, the creator endows its child with initial wealth  $W_{\text{init}}$  and takes a share of the child’s profit. Typically we run with each agent paying one-tenth of its profit plus a small constant sum to its creator, but performance is little affected if this share is anywhere between 0 and .25. Each agent also pays *resource-tax* proportional to the number of instructions it executes. This forces agent evolution to be sensitive to computational cost. Agents are removed if, and only if, their wealth falls below their initial capital, with any remaining wealth returned to their creator. Thus the number of agents in the system varies, with agents remaining as long as they have been profitable.

This structure of payments and capital allocations is based on simple principles[Bau98]. The system is set up so that everything is owned by some agent, interagent transactions are voluntary, and money is conserved in interagent transactions (i.e. what one pays another receives)[MD88]. Under those conditions, if the agents are rational in that they choose to make only profitable transactions, a new agent can earn money only by increasing total payment to the system from the world. But irrational agents are exploited and go broke. In the limit, the only agents that survive are those that collaborate with others to extract money from the world.

Hayek guarantees everything is owned by auctioning the whole world to one agent, which then can alone act on the world, sell it, and receive reward from it. Agent interactions are voluntary in that they respect property rights. For example, the agent owning the world can refuse to sell by outbidding other agents, which costs her nothing since she pays herself. The guide for all *ad hoc* choices, e.g. the one-tenth profit fraction, was that the property holder might reasonably make a similar choice if given the option. In principle, all such constants could be learned.

By contrast such property rights are not enforced in real ecologies[MD88], simulated ecologies like Tierra[Ray91] and Avida[LOCA99] or other multi-agent learning systems such as Lenat’s Eurisko[Len83, MD88] or Holland classifier systems[Hol86]. When not everything is owned, or money is not conserved, or property rights are not enforced, agents can earn money while harming the system, even if other agents maximize their profits. *The overall problem then can not be factored because a local optimum of the system will not be a local optimum of the individual agents.* For example, because in Holland Classifiers many agents are active simultaneously, there is no clear title to reward, which is usually shared by active agents. A Tragedy of the Commons[Har68] can

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<sup>1</sup>This is in the spirit of the currently best-selling AI textbook([RN95]) which defines (p7): “an agent is just something that perceives and acts.”

then ensue in which any agent can profit by being active when reward is paid, even if its action harms performance of the system[Bau96]. Such systems evolve complex behavior, but not accurate credit assignment[MD88, Bau98].

We have done runs with conservation of money broken in various ways. The system learns to exploit any way of “creating money” without solving the hard problems posed by the world, c.f.[Bau96]. Conversely if money leaks out of the system too fast (the resource-tax is a small leak, but must be kept very small, less than  $10^{-5}$  per instruction) the economy collapses to essentially random agents. We have also experimented with violations of property rights. For example, we did runs identical to Hayek in every way, except using the procedure advocated in Zeroth Level Classifier Systems[Wil94] of choosing the winning bidder with probability proportional to its bid. These runs never solved more than 6-block problems, even if given *NumCorrect*, nor made any progress on Rubik’s cube.

### 3 Blocks World

We have trained Hayek by presenting Blocks World problems of gradually increasing size, cf fig. 1. Each problem contains 4 stacks of colored blocks, with with  $2n$  total blocks and  $k$  colors. The leftmost stack, stack 0, serves as a template only and is of height  $n$ . The other three stacks contain, between them, the same multi-set of colored blocks as stack 0. The learner can pick up the top block on any but stack 0 and place the block on top of any stack but 0. The learner takes actions until it asserts “*Done*”, or exceeds  $10n \log_3(k)$  actions. If the learner copies stack 0 to stack 1 and states *Done*, it receives a reward of  $n$ . If it uses  $10n \log_3(k)$  actions or states *Done* without copying stack 0, it terminates activity with no reward. Note the goal is to discover an *algorithm* capable of solving random new instances. The best human-generated algorithm of which we are aware (a recursive algorithm similar to that solving the Tower of Hanoi problem that we developed in conversation with Manfred Warmuth) is capable of solving arbitrary Blocks World problems in  $4n \log_2(k)$  grabs and drops.

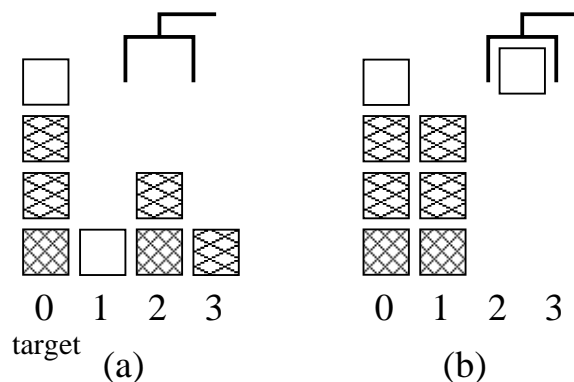


Figure 1: A Blocks World Instance with 4 blocks and 3 colors. (a) shows the initial state. (b) shows the position just before solution. Hatching represents color. When the hand drops the white block on stack 1, the instance will be solved and Hayek will see another random instance.

Blocks World (BW) has been studied in hundreds of papers (cf [Kor87, WB91, Koz92, Koe98, BK96, DBDR98]) because, while easy for humans to solve or program, its exponential-size state-space and the complexity of code that would solve it have stymied fully autonomous methods. One line of research has involved planning programs, which are told the goal and do an extensive search

to attempt to solve an individual instance. Until about 1998, state of the art planning programs such as Prodigy 4.0 could only solve problems involving about 5 blocks before succumbing to the combinatorial explosion[BK96]. More recently, it has been discovered that the Blum-Furst graphplan heuristic[BF97] can be effectively used to constrain the search, postponing although not preventing the combinatorial explosion. The best results are by Koehler[Koe98] who solved a related 80-block stacking problem in 5 minutes and a 100 block problem in 11 minutes. (Note the exponential increase.) Koehler’s problems are simpler than those discussed here in that all blocks start on the table, with no block on top of them, avoiding any necessity for the long, Tower-of-Hanoi-like chains of stacks and unstacks required by a 3-block wide table.

No standard method appears able to address the Blocks World reinforcement learning problem. Previous attempts to reinforcement learn in Blocks World (e.g. [WB91, Koz92, BP94, DBDR98]) addressed only much simpler versions not involving abstract goal discovery (i.e. they attempted to stack a fixed order of colors rather than match an arbitrary target stack); nor did they have a fixed small table size; nor did they involve goal stacks of more than a handful of blocks.

Q-learning[Wat89] is inapplicable to this problem because there are combinatorial numbers of states—there are almost  $10^{100}$  distinct possible 200 block goal stacks and for each of these, a search space of some  $10^{100}$  states, among which there is a single goal state. The standard approach to dealing with large state spaces is to train a neural net evaluator in TD( $\lambda$ ) learning, c.f. [SB98]. This has been effective in some domains, e.g. backgammon[Tes95]. However the success at backgammon was attributed to the fact that a linear function is a good evaluator[Tes95]. TD( $\lambda$ ) using a neural net evaluation function is not suitable here because the goal is so nonlinear and abstract, and the size of the problem description is variable<sup>2</sup>. We nonetheless attempted to train a neural net<sup>3</sup>. This succeeded in solving 4 block problems from a raw representation or 8 block problems using a powerful hand-coded feature *NumCorrect*. *NumCorrect* returns the current number of correct blocks on stack1, i.e. the largest integer  $l$  such that the bottom  $l$  blocks of stacks 0 and 1 agree in color. A recent attempt at inductive logic programming solved only 2 block problems[DBDR98]. We have experimented extensively with Genetic Programming[BD99] and also some less standard approaches: a hill-climbing approach on computer programs, previous economic models, and an attempt to induce an evaluation function using S-expressions. These approaches all solved at most 4 or 5 block problems reliably<sup>4</sup>. See Appendix A for a fuller discussion of these results.

## 4 Representation Language

This section discusses the syntax of the agent programs. The programs of the agents discussed in this paper are typed S-expressions[Mon94]. S-expressions are recursively defined as a symbolic expression, as in Lisp, consisting of either a symbol or a list structure whose components are S-expressions. S-expressions are isomorphic to parse trees. A simple example is shown in figure 4 (a).

All our expressions are typed, either taking integer, void, color, or boolean values. All operations respect types so that colors are never compared to integers, for example. The use of typing semantically constrains mutations and thus improves the likelihood of randomly generating

<sup>2</sup>For a discussion of various other difficulties see [WB91]

<sup>3</sup>We report in this paragraph results on a substantially simpler version of the problem where, rather than demanding that the learner say “done” when the stack is correct, we externally supplied the “done” statement. Run on the full problem, these standard methods did worse.

<sup>4</sup>A simpler economic model learned to solve arbitrary Blocks Worlds problems, but only if provided intermediate reward in training whenever an action was taken partially advancing solution[Bau96].

meaningful and useful expressions.

Our S-expressions are built out of constants, arithmetic functions, conditional tests, loop controls, and four interface functions. *Look(i,j)* returns the color of the block at location  $i, j$ . *Grab(i)* and *Drop(i)* act on stack  $i$ . *Done* ends the instance. Some experiments also contain the function *NumCorrect*, and some contain a random node  $R(i, j)$  which simply returns  $i$  with probability  $1/2$  and  $j$  with probability  $1/2$ .

The system starts with a single, special, hand coded agent, called *Seed*, with zero wealth. *Seed* does not bid, but simply creates children as random S-expressions in the same way that expressions are initiated in Genetic Programming[Koz92, Mon94], choosing each instruction in the S-expression at random from the instruction set with increasing probability of choosing a terminal expression so the tree terminates.  $W_{\text{init}}$  is initially set to 0, so all agents can create freely. Eventually one of *Seed*'s descendants earns reward. Thereafter  $W_{\text{init}}$  is set equal to the size of the largest reward earned, and agents can no longer create unless they earn money.

We report elsewhere experiments using a similar economic model, but with three other representation languages. Hayek1[Bau96] used simple productions of a certain form in some ways similar to the language used by Classifier Systems[Hol86]. In Hayek and also in classifier systems, a configuration is only stable if agents are on average followed by agents of higher bid, else agents will go broke. Conversely, a human writing programs in such a language will naturally use higher bids to prioritize agents, which may result in high bidding agents preceding lower bidding ones. This means that many programs that a human could write in a given language may not be dynamically stable. Although the Classifier System language has been proved universal[For85], it is not clear whether it remains universal when restricted to stable configurations. ([Bau96] and [LU99] independently pointed this problem out.) Similarly, because of this problem, Hayek1 was only able to solve large Blocks World problems when given intermediate reward for partial progress.

Hayek2 used an assembler-like language inspired by Tierra[Ray91]. This proved both ineffective, and incomprehensible to humans.

We call the system experimented with here, using typed S-expressions, Hayek3.

Hayek4, reported on in a companion paper, uses Post Production systems. This language is Turing complete, and Hayek4 succeeds in learning a program capable of solving arbitrary Blocks World problems (of our type). By comparison, we do not believe the S-expression language used in Hayek3 is Turing complete, and so we will see that Hayek3 only succeeds in building systems capable of solving large, but finite, Blocks World problems.

## 5 Blocks World Behaviors

We now report on experiments running the Hayek3 system described in the last section. We trained Hayek on the following distribution of problems:

- One block problems if there was no money in the system, e.g. until the first instance is solved.
- Else we presented problems with size uniformly distributed between 1 and  $2 + m + m/5$ , where  $m$  was set as follows. We maintained a table  $S_p(i)$  as the fraction of the last 100 instances of size  $i$  presented which were solved.  $m$  was increased by 1 whenever  $S_p(m) > 75\%$  and  $S_p(m + 1) > 30\%$ .  $m$  was decreased by 1 whenever  $S_p(m) < 25\%$  and  $S_p(m - 1) < 70\%$ .

This distribution was chosen ad hoc, with a small amount of experimentation, to present larger instances smoothly as Hayek learned. The reason why we used different conditions for increasing and decreasing  $m$  was to prevent rapid oscillation between two neighboring values.

We first describe experiments with *NumCorrect* in the language. Within a few hundred thousand training instances, Hayek3 evolves a set of agents that cooperate to solve 100-block goals virtually 100% of the time and problems involving much higher goal stacks over 90% of the time. See fig. 5 a. Our best run to date solved almost 90% of 200-block goals, the largest stacks among our tests, in a problem with  $k = 3$  colors. Such problems have about  $10^{100}$  states. Hundreds of agents act in sequence, each executing tens of actions, with a different sequence depending on the instance. The system is stable even though reward comes only at the end of such long sequences. The learning runs lasted about a week on a 450MHz Pentium 2 Processor running Linux, but the trained system solved new 200-block problems in a few seconds.

Empirically we find Hayek3 evolves a system with the following strategy. The population contains 1000 or more agents, each of which bids according to a complex S-expression that can be understood, using Maple, to be effectively equal  $A \cdot NumCorrect + B$ , where  $A$  and  $B$  are complex S-expressions that vary across agents but evaluate, approximately, to constants. The agents come in 3 recognizable types. A few, which we call “cleaners”, unstack several blocks from stack 1, stacking them elsewhere, and have a positive constant  $B$ . The vast majority, i.e. about 1000, which we call “stackers”, have similar positive  $A$  values to each other, small or negative  $B$ , and shuffle blocks around on stacks 2 and 3, and stack several blocks on stack 1. “Closers” bid similarly to stackers but with a slightly more positive  $B$ , and say *Done*.

At the beginning of each instance, blocks are stacked randomly. Thus stack 1 contains about  $n/3$  blocks, and one of its lower blocks is incorrect. All agents bid low since *NumCorrect* is small, and a “cleaner” whose  $B$  is positive thus wins the auction and clears some blocks. This repeats for several auctions until the incorrect blocks are cleared. Then a stacker typically wins the next auction. Since there are hundreds of stackers, each exploring a different stacking, usually at least one succeeds in adding correct blocks. Since bids are proportional to *NumCorrect*, the stacker that most increases *NumCorrect* wins the auction. This repeats until all blocks are correctly stacked on stack 1. Then a closer wins, either because of its higher  $B$  or because all other agents act to decrease the number of blocks on stack 1 and thereby reduce *NumCorrect*. The instance ends successfully when this closer says *Done*.

A schematic of this procedure is shown in figure 2.

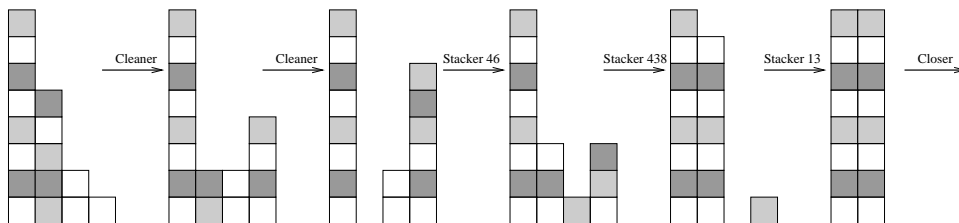


Figure 2: An example of solution of an 8 block instance.

This evolved strategy does not solve all instances. It can fail, for example, when the next block to place is buried under 5 or 10 other blocks, and no agent can find a way to make progress. The next block needed is more likely to be buried deep for higher numbers of colors, so runs with  $k = 10$  or  $k = 20$  colors evolve more agents, up to 3000 in some experiments, and hence run slower than runs with  $k = 3$ . Nonetheless, they follow a similar learning curve and have learned to solve BW instances with up to 50 blocks consistently.

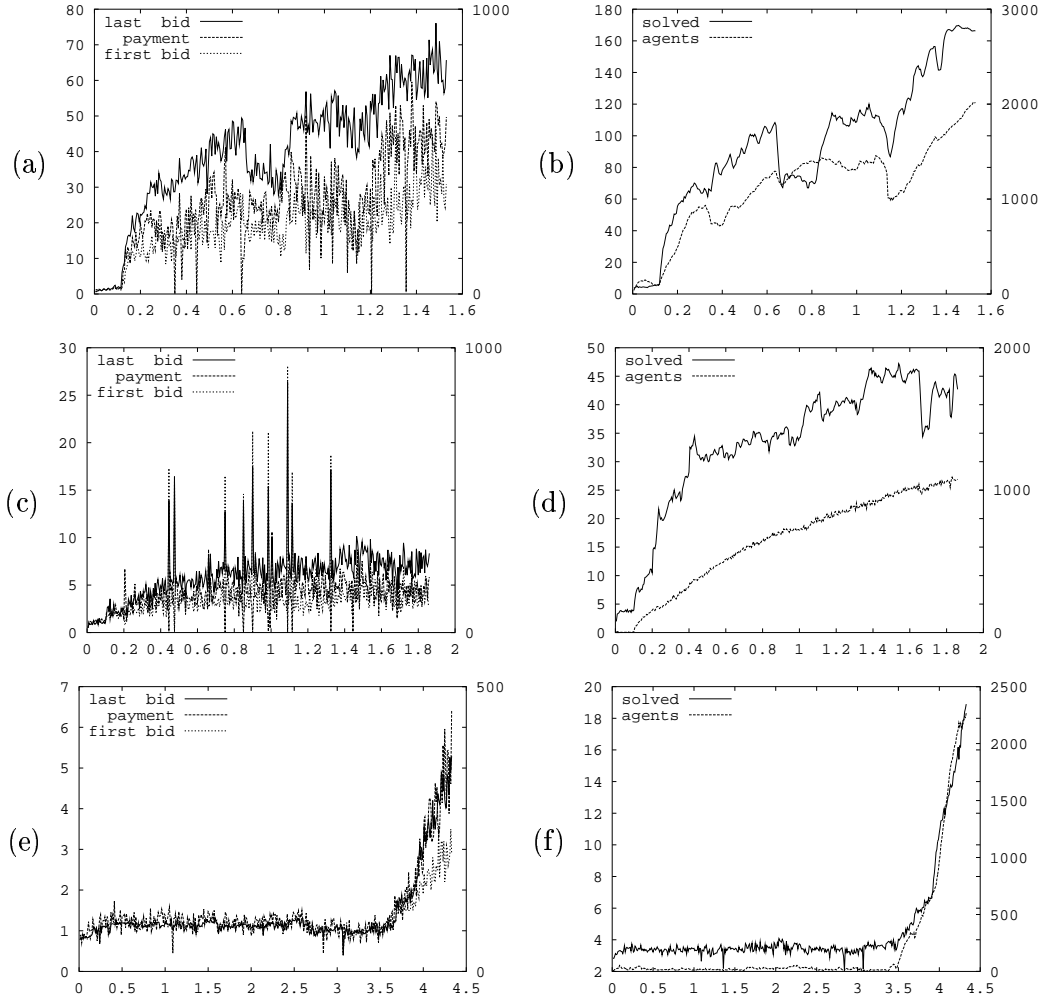


Figure 3: Data from Hayek3 runs (a)(b) with NumCorrect, (c)(d) without NumCorrect or  $R$  nodes, and (e)(f) without NumCorrect but with  $R$  nodes. The horizontal axis is in millions of instances presented. The right vertical axis is in units of agents. The left vertical axis is in units of numbers of blocks or, equivalently, reward. (a), (c), and (e) show the moving average over the last 100 instances of winning first and last bid and payoff from the world. (b) (d) and (f) show, in solid, the moving average of the score, computed as  $\sqrt{2 \sum_{i=1}^{200} p(i) i}$  for  $p(i) =$  fraction instances of size  $i$  solved. If all instances of every size were solved, score would be about 200. Score is computed only over instances with no new agents introduced. This approximates the average size instance being solved if agent creation is suspended. The moving average of payoff is lower than score because new agents frequently lead to failures. A system performing at the level of score can, however, be achieved at any time by suspending creation. (b) (d) and (f) show in dashed line the moving average of the number of agents in the population (measured against the right hand axis). The run shown in (c) and (f) has seen more instances than the others because it took a long time to discover how to solve effectively, hence was being presented very small instances, with very small population. Effective solving coincides with a split of first and last bids, when it learns to estimate value of states.



Classical planning programs search a large space of possible actions. By contrast, Hayek3 learns to search only a few hundred macro-actions and learns to break BW into a sequence of subgoals, long an open problem in the planning community[Kor87]. Note the macro-actions are tuned to the problem domain. For example, the macro-action of the cleaner involves unstacking from 1 and stacking elsewhere, and the macro-action of each stacker involves unstacking some block or blocks and placing it or them on stack 1.

Standard RL approaches, as well as Classifier Systems, try to recognize progress after a single action. Such an evaluation requires not only understanding *NumCorrect*, but also concepts such as “number of blocks on top of the last correct block”, “next block needed”, and “number of blocks on top of the next block needed”. It is too complicated to hope to learn in one piece. Hayek3 instead succeeds by simultaneously learning macro-actions and evaluation. Its macro-actions make sufficient progress that it can be evaluated.

We also ran Hayek3 without the function *NumCorrect*. *NumCorrect* can still be computed as

$$For(And(EQ(Look(0, h), Look(1, h)), Not(EQ(Look(0, h), Empty)))).$$

Hayek3 learned to approximate *NumCorrect* as *For((EQ(Look(0, h), Look(1, h))))* and, employing a similar strategy as above, solved problems involving goal stacks of about 50 blocks, c.f. fig. 5b.

To improve *For((EQ(Look(0, h), Look(1, h))))* to the full *NumCorrect*, Hayek3 has to make a large jump at once. We have discovered that adding a node  $R(a, b)$  to the language greatly improves the evolvability.  $R(a, b)$  simply returns the subexpression  $a$  with probability 1/2 and the subexpression  $b$  with probability 1/2. In addition, we add mutations changing  $R(a, b)$  to  $R(a, a)$  or  $R(b, b)$ . The profitability of an S-expression containing an  $R$  node interpolates between that of the two S-expressions with the  $R$  node simply replaced by each of its two arguments. This seems to smooth the fitness landscape. If one of the two alternatives evolves to a useful expression, the mutation allows it to be selected permanently.

With  $R$  added to the language, Hayek3 consistently succeeds in discovering the exact expression for *NumCorrect*, c.f. fig. 5c. The runs then follow a strategy identical to that with *NumCorrect* included as a primitive. Unfortunately, the runs are slower by a factor of perhaps 100 because the single built-in instruction *NumCorrect* becomes a complex *For* loop with an execution cost in the hundreds of slow instructions. Accordingly, after a week of computation, the system has learned only to solve 20-block problems consistently. Nevertheless, this system is stably following a similar learning curve to that with a primitive *NumCorrect* supplied, apparently learning only a constant factor slower, which could be made up with a faster computer. By contrast, standard strategies typically incur an exponentially increasing cost for solving larger problems. The 50 and 20 block problems solved by Hayek3 without *NumCorrect* are each several *times* bigger than were solved by competing methods with *NumCorrect*.

## 6 Meta-learning

This section describes experiments with a more complex approach where new agents are created by existing agents, rather than as simple mutations of existing agents. The point of this is meta-learning. By assigning credit to agents that create other good agents, we hope to learn how to create good agents, and thus to greatly expedite the learning process.

In this scheme there are two kinds of agents: solvers and creation agents. An agent is a creation agent if the instruction at the root of its S-expression is a “modify” or a “create”, else it is a solver. Solvers behave as the agents we discussed in previous sections.

Creation agents don't bid. Instead, in each auction all the creators that have wealth more than  $W_{init}$  ( a fixed sum) are allowed to create. The creator endows its new child with an initial wealth  $W_{init}$ . Creators are compensated for creating profitable solvers, for creating profitable creators, or for creating agents modified into profitable agents, as follows. In each instance, all agents pay one tenth their profit (if any) plus a small constant sum toward this compensation. For agents created de novo, this payment goes to their creator (which then recursively passes one tenth on to its creator). For agents created by a "modify" command, the payoff is split equally between their creator and the holders of intellectual property used to create them. Intellectual property rights in an agent  $A$  are deemed shared equally by its creator and, recursively, the holder of intellectual property rights in the agent (if any) modified to create  $A$ .

Creators that profit create more agents. Creators with no surviving sons and wealth less than  $W_{init}$  are removed. Creators are thus removed when their wealth, plus that of all surviving descendants (which they are viewed as "owning") is less than  $W_{init}$  so that creators, like solvers, are removed whenever their net impact has been negative. Thus there is evolutionary pressure toward creators that are effective at creating profitable progeny.

The system is initiated with a single creation agent, called Seed, with zero capital., Seed is written (by us) to create children that consist of random code. These agents can in turn create other agents, if their code so indicates. As in the simpler scheme described in section 2,  $W_{init}$  is initiated as zero, but raised when an agent first earns money from the world. Creators endow their child with  $W_{init}$  capital.

In addition to the instruction set described in section 4, Creators employ one "wildcard" \* of each type, and 4 binding symbols \$1,\$2,\$3,\$4 of each type, and two functions *create* and *modify*. These are discussed below.

Wildcards are treated just as symbols of appropriate type when they appear in creators. When a wildcard (or an unbound binding symbol) appears in an S-expression of a solver, however, it is expanded into a random subtree. This is done similarly to the general approach in strongly typed genetic programming[Mon94], by growing from the top down. At each node, one randomly chooses an expression of the correct type. Unless this expression is a constant (i.e. takes zero arguments) we must iterate to the next level, choosing random expressions for the children (subexpressions). At each step, we multiplicatively decrease the probability of generating a non-terminal symbol so that the expansion terminates with a small tree<sup>5</sup>.

*Create* takes one argument. *Create* creates a new agent whose expression is simply identical to its argument.

The *modify* instruction has two arguments. *Modify* chooses a random agent in the population<sup>6</sup>, and attempts to match the expression in its first argument with a subexpression of either the randomly chosen agent's void or integer expressions. If it fails to find a match, no agent is created. If it succeeds, it creates a new agent identical to the matched one except that *modify's* second argument is substituted in place of the matched expression. See figure 4. A match must be exact, except that the arguments of the *modify* expression may contain *binding symbols* and *wildcards*. Binding symbols and wildcards can match arbitrary expressions of the appropriate type. The

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<sup>5</sup>Let  $P_{expand}(d)$  be the probability of choosing a non-terminal at depth  $d$ . We set  $P_{expand}(d+1) = P_{expand}(d)*C$ . We imposed limit conditions that  $P_{expand}(\text{depth} = 0) = \text{initial-expansion}$  and  $P_{expand}(\text{average-depth})=0.5$ . We chose (ad hoc, without experimentation) an average-depth of 3 and initial-expansion of 0.9. The constant  $C$  is then computed as:  $C = \exp((\log(0.5) - \log(\text{initial-expansion}))/\text{average-depth})$ .

<sup>6</sup>Note that there is no "selection" operator per se. *Modify* chooses a random agent in the population. Selection occurs because unprofitable agents are removed from the population. Modifiers can only sense which agent they modify in that they must match a pattern. We have not yet explored allowing modifiers other sensations and options, such as choosing to modify a wealthy agent, or even one recently active.

difference between them is that such binding symbols in the first argument are bound when it matches, and if the same binding symbol also occurs in the second argument, the binding symbol is replaced by the expression it was bound to. Wildcards match independently and do not bind. Every tenth new agent, however created, is further mutated by having a randomly chosen node or nodes replaced by random trees. This helps to ensure the system will not get stuck in a configuration from which it can not further evolve.

Note that a modify operation may, e.g., use code from a solver in the population to create another creator. This creator may, e.g., then use code from a second solver in creating another solver. In this way it is possible for code fragments from different agents to be combined.

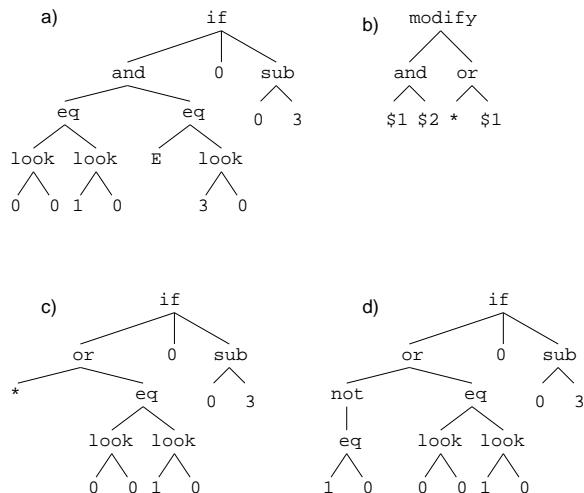


Figure 4: Figure (a) shows an integer expression. This is equivalent to  $if[(look(0,0)=look(1,0))and(look(3,0)=E)] then 0 else -3$ . It returns 0 if the bottom block of the 0 stack is the same color as the bottom block of the 1 stack and there are no blocks on the 3 stack, else it returns  $-3$ . Figure (b) shows the void expression of a creator, with *modify* at its root. This can match with the expression of figure (a) by matching the *and* expressions, binding \$1 to  $eq(look(0,0),look(1,0))$  and binding \$2 to  $eq(E,look(3,0))$ . Substituting the second argument of the modify (namely the subtree rooted by *or*) results in Figure (c). Finally the wildcard  $*$  is expanded into a random tree, resulting in the expression of Figure (d).

When we ran on Blocks World with creation agents, the system learned a collection of solvers employing a similar strategy to that described in section 5. The behavior of *creation* agents was less clear. Some evolved to modify sequences of grabs and drops in ways that create useful stackers or cleaners. Others were not transparent to us.

We ran control experiments to discover whether we were in fact able to learn effective ways to create. We first compared a run with the creation agents to a control in which creation agents attempted to create, but we discarded the agent they created in favor of a mutated agent. Hayek3 with creation agents performed 30-50% better after a week of execution than this control. This indicates that the creation agents were learning useful techniques. However, the pattern matching in the creation agents was expensive in time, and the system with creation agents performed no better than the less complicated system reported in previous sections.

The conclusion of our experience with meta-learning is thus mixed. With our current techniques, meta-learning is probably not a practical alternative. It remains plausible that with a more powerful creation language, and for yet more complex problem domains, meta-learning may reemerge as a

useful technique.

## 7 Rubik’s Cube

We tried to learn a Hayek capable of unscrambling a randomly scrambled Rubik’s cube. This section will assume some familiarity with Rubik’s cube. For the uninitiated: a “cubie” is one of the 26 physical little cubes out of which a Rubik’s cube is constructed. We call a “square” one face of a cubie.

Rubik’s cube is a complex problem. We experimented with a variety of instruction sets, including several versions of *numcorrect* functions, several presentation schemes of increasingly harder instances, and several reward schemes. All of the versions we experimented with used creation and modify operators, as discussed in the previous section.

We used instruction sets with several types. For example, one set of runs used 6 types: boolean, integer, coordinate, face, square, and void (i.e. actions). To describe a particular square you could write  $s(i, x, y)$  where  $i$  is a coordinate of type “face” specifying a face of the cube (e.g. the face with red center), and  $x$  and  $y$  are coordinates taking values 0,1, or 2, specifying the location on the face. Having specified two squares  $s_1$  and  $s_2$ , it would be possible to compare their color with the boolean-valued function  $EQCs(s_1, s_2)$ , which returns true if and only if they have the same color. We supplied an instruction *Sum* taking a Boolean argument, which summed any constructed boolean function over all cube faces. Hayek evolved evaluation functions which used sums to estimate the number of correct cubies. We realized, however, that it was impossible for Hayek to express the physical concept of a “cubie” in this language, and so also experimented with a language based around cubies, in which we supplied a function *NumCorrect* describing the number of correct cubies.

We also worked with two different presentation schemes. In *Presentation Scheme 1*, we initially presented cubes scrambled with one rotation, and as Hayek learned to master cubes at a given level of scrambling, presented cubes scrambled with increasing numbers of rotations. We then gave reward only for completely unscrambling the cube. In *Presentation Scheme 2*, we presented a completely scrambled cube, i.e. a cube scrambled with 100 random rotations, and presented Hayek with reward proportional to its partial progress at the time it said “done”. For this purpose we measured partial progress according to three different metrics (in different runs): (a) the number of cubies it got right, (b) the number of cubies it got right according to a fixed sequence (i.e. we numbered the cubies starting with the top face and gave Hayek reward equal to one less than the first cubie in the sequence that was wrong), and (c) the number of cube faces it got right (maximum 54).

We also worked with two different cube models. *Cube model A* allowed three actions: F– a one quarter move on the front face, and X and Y rotations of the whole cube. *Cube model B* held the X, Y, and Z axes of the cube fixed, but allowed a one-quarter move on each of six faces (L,R,F,B,T,D).

Some representative results are as follows. In a run using Presentation Scheme 1, with Hayek given an instruction *NumCorrect* that we set equal the total number of cubie-faces correct relative to the center cubie-face, Hayek learned after a few days to descramble cubes scrambled with 20 quarter-moves 80% of the time and cubes scrambled with 30 quarter-moves 40% of the time.

In another experiment, using Presentation Scheme 2, with reward equal to the total number of cubie faces it left correct, Hayek learned to improve scrambled positions by about 20 cubie faces. (I.e. the average number of correct cubie faces at the end minus number of correct cubie faces at the beginning was about 20.)

In runs with Presentation Scheme 2, with Hayek given reward for partial success defined as

number of correct cubies in a particular (x,y,z) order, stopping at the first incorrect cubie (and not counting the center cubies, which are always correct by definition), Hayek learned to solve about ten cubies or the whole first slice and two more cubies of the middle slice. This is further than many humans get.

The bottom line is that for each of the schemes we tried, Hayek made significant progress, but after a few days of crunching it would plateau and cease further improvement. Actually, this describes as well the progress of many humans, for the simple reason that, as the cube becomes more solved, it becomes increasingly difficult to find operators which will make progress without destroying the progress already achieved. Such operators become increasingly long chunks of code, and there are no evident incremental properties of the chunks, so it becomes exponentially hard to find them unless some deep understanding can be achieved.

Hayek’s strategy in these runs was fundamentally similar to its strategy in Blocks World: it produced a collection of three to seven hundred agents that took various clever sequences of actions and bid an estimate of the value of the state they reached. Typical instances involved a dozen or so auctions, the winner in each making partial progress. Hayek made progress in all of the different schemes, and it was not obvious that any scheme was particularly preferable to the others.

It is unclear whether Hayek, at least using the S-expression representation here, can exploit compact structure of the problem domain in more powerful ways. It is worth noting, for example, that within the S-expression languages used here, it does not seem that high level concepts such as “cubies” can be represented (never mind learned) to the extent that we do not simply insert them by hand. However, it also seems likely there is fundamentally less structure to be exploited in Rubik’s cube than in Blocks World. The smallest program a human could write that will solve Rubik’s cube is much much larger than the smallest human-crafted program to solve Blocks World, while at the same time the state space in Blocks World is huge (in fact infinite) and the state space in Rubik is merely large. Indeed Rubik can be solved by brute search[Kor97].

## 8 Rushhour

This section describes experiments applying Hayek to the commercially available game “Rush Hour”. The game comes with 40 different puzzle instances. Figure 5 shows instance 35. Each instance consists of several cars placed on a 6x6 grid. Cars are of width 1 and come in two different lengths, 2 or 3. A car can only be moved forward or backward, not turned. The goal of the game is to move the cross marked car out the exit. The game was honored for excellence by “Mensa” and humans find the problems challenging. Moreover the game is provably hard: [FB99] shows that the generalization to an  $n \times n$  grid is P-space complete, equivalent to a reversible Turing machine, and that solution may require a number of moves exponential in  $n$ .

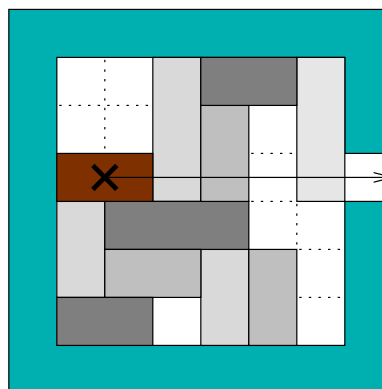


Figure 5: Rush hour problem 35

We trained Hayek, and also a Genetic Program, by presenting problems of gradually increasing difficulty, including easy training problems generated from the original 40 by randomly removing cars. Rush Hour presents new challenges particularly in the selection of a representation language for the agents, since agents need flexibility in specifying actions and must be able naturally to express relations such as “the car that is blocking this

car.” The version of the program discussed here used integer-type S-expressions, and used creation operators. We call this Hayek3.1

Agents were limited to expressions of up to 20,000 nodes. The language consists of constants (0,1,2), arithmetic functions (add, subtract, multiply), if-then-else, and some additional functions as follows. There is a pointer which is initiated pointing to the marked car. The function *PUSH* moves the car pointer to either the car blocking forward motion of the pointed car or the car blocking its backward motion, depending on its whether its argument is zero or not. (If there is no such blocking car, i.e. motion in that direction is blocked only by a wall or alternatively it is possible to drive off the board, then *PUSH* fails and takes no action.) As a side effect, the previously pointed car is pushed onto a car stack. Cars can be popped back giving our language a weak backtracking ability. Function *PUSH2* behaves like *PUSH*, except that it moves the car pointer to the second blocking car (i.e. the car that would block next if the first blocking car were removed). Function *MOVE* moves the pointed car forwards or backwards. Each of these functions returns 1 if they succeed and 0 if they fail.

Creators employ additional statements: *MODIFY*, a “wildcard” \*, and 4 binding symbols \$1,\$2,\$3,\$4.

To get a rough estimate of the difficulty of the problems, we applied random search to them. Table 8 shows the problems ordered by the number of moves used to solve them by a random search algorithm. We call the ranking according to table 8 “order no.” and the number supplied by the manufacturer “task no.”.

Order	<b>1</b>	<b>2</b>	3	<b>4</b>	5	<b>6</b>	<b>7</b>	<b>8</b>	<b>9</b>	<b>10</b>
task:	<b>9</b>	<b>13</b>	28	<b>14</b>	23	<b>10</b>	<b>7</b>	22	<b>2</b>	<b>12</b>
moves:	<b>1394</b>	<b>3181</b>	5909	<b>6063</b>	6069	<b>6854</b>	<b>7487</b>	8931	<b>9919</b>	<b>10204</b>
Order	<b>11</b>	<b>12</b>	<b>13</b>	14	15	<b>16</b>	<b>17</b>	18	19	20
task:	<b>4</b>	<b>6</b>	<b>20</b>	27	16	<b>1</b>	<b>21</b>	25	30	17
moves:	<b>11718</b>	<b>13134</b>	<b>13532</b>	14281	14451	<b>20916</b>	<b>21968</b>	22348	23255	24336
Order	21	<b>22</b>	<b>23</b>	24	25	26	27	28	29	<b>30</b>
task:	32	<b>3</b>	<b>5</b>	38	37	11	15	19	29	<b>8</b>
moves:	25813	<b>26077</b>	<b>27572</b>	27828	28127	36892	41421	57901	58439	<b>70804</b>
Order	31	32	33	34	35	36	37	38	39	40
task:	18	26	24	40	39	34	31	33	36	35
moves:	79533	87473	100016	111243	123884	157495	160507	171687	184341	209923

Table 1: Ordering of the tasks by random-moves. Task is the number supplied by the manufacturer. Moves is the number of moves taken to solve the problem using random search. Problems printed in bold were solved by Hayek3.1.

For each task we allowed a total number of car moves bounded by  $100 + \#random\text{-moves}/100$  where  $\#random\text{-moves}$  was the average number of moves taken to solve that particular instance random search. In each instance we allowed  $10 + \#random\text{-moves}/1000$  auctions. If the move limit or the auction limit is exceeded the instance ends with no reward.

Hayek3.1 was trained using random problems drawn from a distribution crafted to present easier problems first as follows. We first select a problem randomly and uniformly from the 40 supplied with the commercial game. The instance is then subsampled by removing cars as follows. For each problem number we maintain a count of the fraction of instances of that problem number

solved out of the last 100 times this instance was presented<sup>7</sup>. We remove a number of cars so that the fraction of cars remaining, called  $p$ , is equal to the fraction of the last 100 problems of that number solved (up to round-off, since we must present an integer number of cars). This allows Hayek3.1 to learn from simpler examples, and as it learns on each example, we gradually ramp up  $p$  until we are presenting the full example. The reward  $R$  given for solving an instance was  $R = (\text{orderno.})e^{(4.06p-4.06)}$ . Thus more reward was given for harder instances, and reward is exponentially increased for solving instances with no subsampling. The coefficients in the exponent were determined by fitting  $e^{(ap+b)}$  to be equal to 0.01 for  $p = 0$  and 1.0 for  $p = 1$ .

After training, Hayek contains a collection of hundreds of agents that collaborate to solve some of the problems. Figure 6 shows a sequence of agents acting to solve the task no. 20, (a.k.a. order no. 13). Table 8 shows the sequence of bids and number of actions taken. Notice that the agents recognize as they close on a solution, and bid higher. Note that the final bid is quite accurate (although in this case a slight overbid). (This agent was profiting on other instances, but lost money on this one.) In instances where the problem is ultimately not solved, bids stay low, as agents realize they are not close to solution. The ability of agents to recognize distance from the solution is critical to Hayek’s performance. In each of a series of auctions it chooses the agent that (in its own estimation) makes the most progress eventually converging on a solution. This breaks down a lengthy search into a sequence of operations effectively achieving subgoals.

auction #	winning agent	winning bid	Reward	actions
0	a1052	6.749	0	14
1	a1011	7.369	0	15
2	a1006	9.208	0	20
3	a1198	13.741	0	22
4	a1198	13.853	13	35

Table 2: Auctions for a Solution of Problem 20.

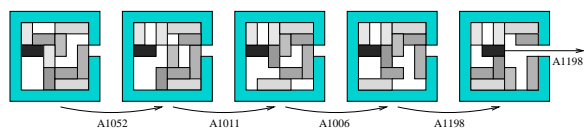


Figure 6: Sequence of positions solving problem 20.

Figures 7 and 8 show the time evolution of statistics for this typical Hayek3.1 run. The data-points are a sample every 10000 instances of the average over 100 instances. Figure 7 shows the winning bid in the first auction, the winning bid in the last auction, and the final payoff. Note that the last bid accurately estimates payoff. Note that first bid is lower.

Figure 8 (a) shows the reward earned. Figure 8 (b) shows the linearly weighted sum of fully solved problems

In each run, Hayek has learned to solve a number of unsubsampled problems. A typical run creates a collection of agents that is solving about 5 unsubsampled problems at a time. The best run solved 9 unsubsampled problems simultaneously. In one run or another Hayek has solved problems 1,2,3,4,5,6,7,8,9,10,12,13,14,20,21 according to task number, i.e. the manufacturer’s ordering. When

<sup>7</sup>Actually, the bins only record data from runs where a new agent did not win the bidding and immediately die. Such failures are deemed irrelevant for purpose of determining the distribution, as this dead agent is no longer part of Hayek3.1 and thus does not affect future performance.

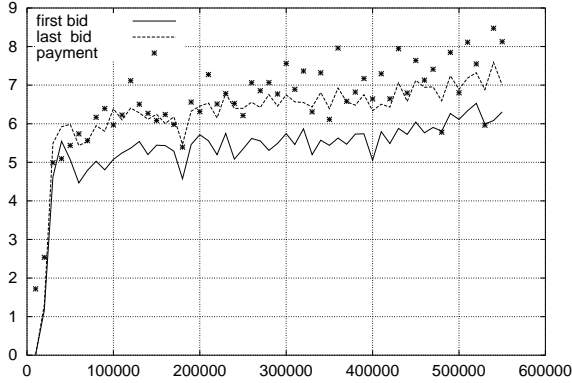


Figure 7: Bid as discriminator between positions

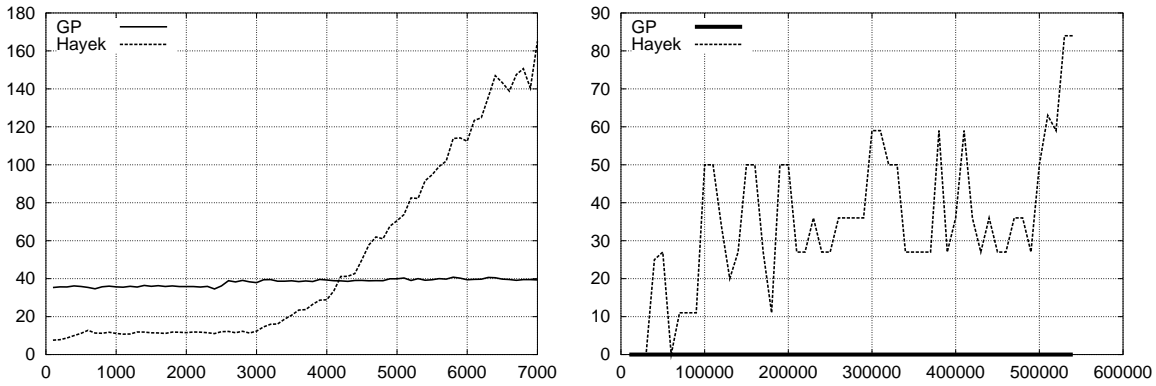


Figure 8: (a) Reward per instance. Horizontal axis times 100 is instances presented to Hayek. Horizontal axis times 10 is GP generations. The first point shown is after 100 GP generations. GP at generation 1 scores about 2 on this scale and learns rapidly at first. Each GP generation involves presentation of 100 instances of each type, i.e. 4000 instances total, to each program in the population. The scales were chosen to roughly agree in wall clock time. GP shows reward per instance of the best program in the population. (b) Score as linearly weighted sum of fully solved instances.

Hayek solves an unsubsampled problem, it generally involves a collaboration of 2 to 5 agents. Subsampled problems are much easier, of course, and are often solved by a single agent.

We have tried several approaches to getting a Genetic Program to solve these problems, varying the instance presentation scheme and other parameters. In each case our GP used the same instruction set as Hayek3.1. The runs shown used a population size of 100 and an instance presentation scheme as similar to that Hayek saw as possible<sup>8</sup>. As the graphs show, while GP learns to solve subsampled problems, and thus to earn reward, it has never learned to solve any of the original problem set. It reaches a plateau where it is solving only subsampled problems after less than one hundred generations. Even when trained solely on Problem 1, we have not been able to learn a program solving the full problem. Possibly some better scheme for ramping up problem difficulty

<sup>8</sup>For a GP, there is a question as to how to compute the fraction of problems of type  $i$  solved in recent generations. The runs shown used the average over the last 10 generations of the fraction of problems of that type solved by the upper half scoring members of the population, which worked as well as any of the other approaches we tried.



would improve performance.

## 9 Discussion

We have created an artificial economy that assigns credit to useful agents. When we enforce property rights to everything and conservation of money, and when agents are rational in the sense that they only accept transactions on which they profit, new agents can enter if and only if they improve performance of the system. Agents are initiated far from rational, but we hoped exploitation of irrational agents would lead to evolution of rational agents and a smoothly functioning system that divides and conquers hard problems. Conversely, we argued that when property rights or conservation of money are not enforced, agents will be able to profit at the expense of the system. Under those conditions, a local optimum of the system will not be evolutionarily stable, and there is no reason to expect evolution to produce a useful system.

We have now performed experiments on Blocks World using three different representation languages for the agents (one in [Bau96], one here, and a very different one reported in [BD00], which adopts a radically different solution strategy), and also with and without creation agents. We have performed experiments on Rubik's cube using two different representation languages (here and in [BD00]). We have reported experiments here on Rushhour. In each of these cases we have stably evolved systems with agents collaborating to solve hard problems. In the Blocks World examples in particular, we are evolving systems with stable sequences of agents hundreds long, receiving reward only at the end. This contrasts with classifier systems, which are rarely able to achieve stable chains more than a few agents long [WG89]. Our dynamics have also been largely robust to parameter variations.

All of this dynamic stability vanishes when conservation of money is broken. If the tax rate is raised too high, so that money leaks out, the system dies. If on the other hand money is non-conserved in a positive fashion, for example by introducing new agents with externally supplied money [Bau96], or in other ways as has happened more than once due to bugs in the code, the system immediately learns to exploit the ability to create money, and cooperation in solving problems from the world disappears. If property rights are broken, for example by choosing the winner of the auction probabilistically as sometimes suggested in the classifier system literature [Wil94], again cooperation immediately breaks and the system can no longer function.

We believe that all of this provides powerful evidence for our economic picture. As we have discussed elsewhere [Bau98] we believe that many of the dynamic phenomena in natural complex systems such as ecologies, as well as many problems with other learning programs, can be profitably viewed as manifestations of violations of property rights or conservation of money. We also propose that Hayek is of interest as an Artificial Life. In particular, our experiments (particularly with meta-learning) answer the open question [TR97] of how to get a collection of self-reproducing programs to cooperate like a multi-cellular being in solving problems of interacting with an external world.

We further conjecture that progress in reinforcement learning in ultra-complex environments requires learning to exploit the compact structure of the problem. This means that learning in complex environments will essentially involve automatic programming. We have here learned to exploit the structure from small problems and generalized the knowledge to produce programs solving large problems.

The problem with automatic programming methods, including ours, is that the space of programs is enormous and has a very bumpy fitness landscape, making it inherently difficult to search. Our proposal to automatically divide and conquer using an economy helps, but the problem remains. This raises the key question of what language will be evolvable.

We have here made progress using S-expressions, which have previously been used in genetic programming[Koz92] precisely because, intuitively, their structure makes mutations more likely to create semantically meaningful expressions. Typing[Mon94] has also been critical here in cutting down the search space.

We found that using a certain generic random node smoothed the fitness landscape in Blocks World and increased evolvability of complex expressions. We found that Hayek was able to solve problems so long as the longest “chunk” of code it needed as a component was not too long to find. Hayek has been able to create very complex programs made of relatively small chunks of code.

Nonetheless, our ability to evolve complex code has been limited. For example, we were not able to progress on Rubik’s cube because after a point the next chunk of code needed was too long to evolve. Meta-learning was proposed as a possible solution to this problem, and our economy was able to support some degree of meta-learning, but to date we have not been able to extend the overall abilities of Hayek using meta-learning.

A related language question is the power of the language. We believe the S-expression language used here is not computationally universal, and accordingly Hayek has not been able to evolve a program solving, for example, arbitrary Blocks World problems. By contrast, we report elsewhere[BD00] a Hayek employing Post Production System language, which is computationally universal, and which succeeds in evolving a simple universal solver for Blocks World. However, this system also founders in the search problem for Rubik’s cube.

More generally, for truly complex problems, one might imagine somehow evolving an information economy, with agents buying and selling information from one another. We have not solved the problem of how to make this work. Also, it would be interesting to have a single, universal language able to address many problem domains, and see if it is possible to transfer knowledge from one domain (e.g. Blocks World) to another (e.g. Rubik’s Cube.)

## 10 Appendix A: Other Methods Applied to Blocks World

For comparison, we tested TD( $\lambda$ ) using a neural net evaluator. For a description of the method, see [SB98] or [Tes95]. We tested several values of  $\lambda$ , several discount rates<sup>9</sup>, and several neural net topologies. We report results on only three color problems. As will be evident, ten color problems would have further stressed the encoding requiring much larger nets if a unary coding were desired. We also report results only on an easier version of our Blocks World where, rather than requiring the learner to recognize when it is *done*, we externally supply that information. We considered a single move to be a grab and drop combination, which improved performance.

Neural nets have a fixed input dimension, so it is unclear how to encode a problem like Blocks World with varying size. An evaluator that does not see the whole space, however, will cause aliasing, turning a Markov problem into a non-Markovian one causing great difficulties. For a discussion of such problems, see [WB91]. After some experimentation, we used neural nets with inputs for each of the top 5 blocks in each stack. We tried both unary and analog encoding for the block color. In unary encoding, each block was represented by 3 inputs to the net: with one having value 1 and the others value 0 to indicate which color the block had and all value 0 indicating no such block (because the stack was not 5 blocks high). In analog encoding, the situation was indicated by a single integer input value from 0 to 3. We ran experiments both with and without *NumCorrect* supplied as an input.

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<sup>9</sup>A discount rate of 1, i.e. no discounting, as used by Tesauro in his backgammon program [Tes95] is the most logical in an episodic presentation scheme like ours. However, empirically discount rates of about 0.9 stabilized learning in some runs.

We presented increasing larger instances as the system learned, as in our experiments with Hayek3. However, to keep learning stable, we found we had greatly to slow the increase in instance size. We presented instances chosen uniformly over the sizes from 1 up to  $l + 1$ , where  $l$  was the largest size being solved 80% of the time.

In experiments without *NumCorrect*, using a 61-20-1 feedforward net, with 60 inputs giving unary encoding of the colors of the top five blocks in each of 4 stacks and one “on” input to realize a threshold, we were able to learn stably and solve problems with  $n = 4$  after several days of learning. No other system without *NumCorrect* supplied did better.

In experiments with *NumCorrect* supplied as an input, the system learned fairly rapidly to solve problems up to about  $n = 8$ . Once it got up to this size, its learning would destabilize and it would crash to much worse performance. Then it would rapidly learn again. The system was never able to learn beyond  $n = 8$ .

These results are perhaps disappointing but not surprising. While neural nets are good at associative memory and approximating some functions, there is no evidence to indicate that they are suited to solving symbolic search problems or to inducing programs as necessary for solving Blocks World. Tesauro’s success with backgammon, for example, is attributable to the fact that a linear function provides a reasonable evaluator for backgammon[Tes95].

We also extensively tested Genetic programming. Results are reported in [BD99]. Since this publication, we have done further extensive tests using population sizes up to 1000, various selection methods, and various combinations of instruction sets including *NumCorrect*. As discussed in [BD99], our results were of some independent interest as bearing on the hotly debated question within the genetic programming community, whether genetic programming benefits from crossover, or would be just as effective using the “headless chicken macromutation” where, rather than swapping random subtrees between trees in the population, random subtrees are simply replaced with new random subtrees. A recent textbook[BNKF98] reports several studies showing these methods essentially equivalent and none showing a big advantage for crossover. Blocks World is thus in some sense an unusually suitable problem for Genetic Programming, as our Blocks World results represent the only published comparison of which we are aware showing crossover much superior to headless chicken. Nonetheless, Genetic Programming had only limited success on this problem. Run with a learner having to say *done*, our best Genetic Programming run solved 70% of size three and 30% of size 4. Run with *done* externally supplied, GP produced at best S-expressions capable of solving 80% of size 4 problems and 30% of size 5. These results are arguably directly comparable to our Hayek results as both GP and Hayek used the same language for their S-expressions.

We also extensively tested hillclimbing approaches using an assembler-like language. Results, reported in [BD99] were that only about 4 block instances could be solved.

We also tested a system which searched for an evaluator as the maximum of a set of S-expressions in Hayek’s language. This system differed from Hayek3 in that we did not learn agents capable of sequences of actions. Rather, as in TD( $\lambda$ ), we considered all pairs of a grab followed by a drop, evaluated the resulting state, and chose the move leading to the highest evaluated state. We modified our evaluator to better estimate the value of the next state reached. There is no standard method for training complex functional expressions as evaluators. We tried to approximate the evaluation from below, by using the max of a number of S-expressions, and removing an S-expression from our population when it overestimated the evaluation, replacing it with another S-expression, generated either randomly or as a mutation of an existing S-expression. This method was also ineffective, solving only single digit blocks world problems. We believe that it is critical to learn macro-actions, as it will be very difficult to learn any evaluator capable of showing progress after only a single grab-drop pair.

[Bau99] discusses several other ineffective approaches we experimented with.

## 11 Appendix B: PseudoCode for Hayek3

```
hayek() {
  for ( ;; ) {
    task.new() // create an instance
    solve()    // try to solve it
    payment()  // perform money transactions
    tax()      // collect taxes
  }
}

solve() {
  record.clean()
  for ( Agent A = Solvers.begin(); A < Solvers.end(); A++ ) A.old_wealth = A.wealth
  while ( task.state() == RUN ) {
    create() // creates new agents
    auction() // performs an auction
  }
}

create() {
  for ( Agent A = Creators.begin(); A < Creators.end(); A++ )
    if ( A.wealth > InitialWealth() ) A.create()
}

auction() {
  best_bid   = -1
  best_state = task
  best_agent = nil
  for ( Agent A = Solvers.begin(); A < Solvers.end(); A++ ) {
    if ( A.wealth > best_bid ) {
      A.move( task, state );
      bid = min( A.eval( state ), A.wealth );
      if ( bid > best_bid ) {
        best_bid   = bid;
        best_state = state;
        best_agent = A;
      }
    }
  }
  if ( best_agent != nil ) { // the winner takes the world into a new state
    task = best_state;
    update_agent_wealth( best_agent, best_bid )
    record.add( best_agent, best_bid, task.reward() )
  }
}
```

```

}

update_agent_wealth( agent, bid ) {
    best_agent.wealth -= best_bid
    if ( record.last = nil ) {
        pay_to_world += best_bid
    } else {
        record.last.wealth += best_bid
    }
}

payment() {
    for ( AuctionRecord R = record.begin() ; R < record.end(); R++ ) {
        delta = R.agent.wealth - R.agent.old_welth
        if ( delta > 0 ) {
            R.agent.wealth += 0.9 * delta
            copyright_payment( A, 0.1 * delta )
        } else {
            R.agent.wealth += delta
        }
    }
}

copyright_payment( Agent A, Money M ) {
    if ( A.father == nil ) {
        A.wealth += M
    } else {
        A.father += 0.5 * M
        copyright_payment( A.father, 0.5 * M )
    }
}

tax() {
    for ( Agent A = Agents.begin(); A < Agents.end(); A++ ) {
        A.wealth -= A.executed_instructions * 1E-6
        A.executed_instructions = 0
        if ( A.wealth < A.initial_wealth AND A.no_living_sons == 0 ) {
            Agents.remove( A )
        }
    }
}

```

## 12 Appendix C: Parameters

There are a number of constants in this system chosen in an ad hoc way, such as the one-tenth profit sharing, the resource-tax,  $W_{init}$  and in versions with creation agents, the equal split between creators and IP holders. These parameters were not subject to evolution within the system. Only a

small amount of empirical tuning was attempted by us. Within reasonable bands in each parameter, the behavior of the system was not found empirically to be qualitatively sensitive to these choices, and we did not establish quantitative differences. Runs of the system require a day or more to learn, and there is considerable random variation from run to run, so empirical tuning is quite difficult unless a qualitatively different behaviour can be observed. We discuss our choices of these values in turn.

Since creators are intuitively considered to own their children, in principle the creator might be allowed to set the profit-sharing fraction as it pleases, say by writing appropriate code in the child, and indeed this was done in an earlier version[BD99]. For the work reported here, just for simplicity, we fixed the profit sharing fraction at one-tenth. We also experimented with one-half. Our subjective impression was that passing one-half of the profit passed too much money to the creation agents, causing an increase in creation which slowed the system. However, the increase, if any, was not huge and we have no statistically significant claim that one-tenth is better than one-half.

$W_{\text{init}}$  was initially set equal to zero. Once there was money in the system because instances were solved,  $W_{\text{init}}$  was simply set equal to the maximum payoff possible in the run. Thus if we were going to run up to a maximum instance size of 200, we set  $W_{\text{init}}$  to 200. This ensured that the created agent would have enough money to make any possible rational bid. If the newly created agent's money ever went below its initial capital, it was removed with its remaining money transferred back to its creator. Thus a creator that invested 200 in a child would not lose the full 200 if the child proved to be unprofitable. We did not experiment with any other choices.

We experimented with several values of the resource-tax. Resource-tax above  $10^{-4}$  per instruction sucked too much money out of the system, which then could not get started. We did not establish a difference between resource-taxes of  $10^{-5}$ ,  $10^{-6}$ , and  $10^{-7}$  although subjectively runs of  $10^{-7}$  ran slower. Our best runs used simply  $10^{-6}$ .

The runs here with creation agents used an equal split of the passed profit between creators and IP holders, but we also experimented with not using IP at all<sup>10</sup>. The equal split of passed profit with IP holders subjectively worked marginally better, but again we do not claim to have established a statistically significant difference.

A more principled approach to all of these choices is a subject for future work. Fixing these parameters by fiat has the feel of fixing prices, which typically can cause inefficiencies in economies.

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<sup>10</sup>Although popular opinion seems to believe that enforcing intellectual property rights are important to progress in real economies, the scientific literature on this subject is wide open, c.f. [Bre76, Mac58].

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